

20875 Software Engineering

Tutorial 2 – errata

- Question 3. The simplest way to solve the exercise is to look for strings that could be a password in the executable file itself. As a (much harder but) more systematic alternative, we can analyse the disassembly of its code. The output of `objdump -M intel -d secret` gives us the following for the `main` function:

```
0000000000401240 <main>:
401240: 41 55          push   r13
401242: 41 b8 2c 30 40 00    mov    r8d,0x40302c
401248: b9 60 30 40 00    mov    ecx,0x403060
40124d: 31 c0          xor    eax,eax
40124f: 41 54          push   r12
401251: ba 90 30 40 00    mov    edx,0x403090
401256: 55              push   rbp
401257: 89 fd          mov    ebp,edi
401259: 53              push   rbx
40125a: 48 89 f3          mov    rbx,rsi
40125d: be 10 30 40 00    mov    esi,0x403010
401262: 48 81 ec 48 18 00 00    sub    rsp,0x1848
401269: 48 8b 3d 90 47 00 00    mov    rdi,QWORD PTR [rip+0x4790]      # 405a00 <stderr@GLIBC_2.2.5>
401270: e8 6b fe ff ff    call   4010e0 <fprintf@plt>
401275: ba 10 00 00 00    mov    edx,0x10
40127a: be 01 00 00 00    mov    esi,0x1
40127f: 48 8b 0d 7a 47 00 00    mov    rcx,QWORD PTR [rip+0x477a]      # 405a00 <stderr@GLIBC_2.2.5>
401286: bf 1c 30 40 00    mov    edi,0x40301c
40128b: e8 40 ff ff ff    call   4011d0 <fwrite@plt>
401290: 31 c0          xor    eax,eax
401292: 48 89 e6          mov    rsi,rsp
401295: bf 2d 30 40 00    mov    edi,0x40302d
40129a: e8 11 ff ff ff    call   4011b0 <_isoc99_scanf@plt>
40129f: 48 b8 4d 79 53 65 63    movabs rax,0x746572636553794d
4012a6: 72 65 74          xor    rax,QWORD PTR [rsp]
4012a9: 48 33 04 24          xor    rax,QWORD PTR [rsp]
4012ad: 48 ba 50 61 73 73 77    movabs rdx,0x64726f7773736150
4012b4: 6f 72 64          xor    rdx,QWORD PTR [rsp+0x8]
4012b7: 48 33 54 24 08    xor    rdx,QWORD PTR [rsp+0x8]
4012bc: 48 09 d0          or     rax,rdx
4012bf: 74 2e          je    4012ef <main+0xaf>
4012c1: 48 89 e6          mov    rsi,rsp
4012c4: bf 43 30 40 00    mov    edi,0x403043
4012c9: 31 c0          xor    eax,eax
4012cb: e8 c0 fd ff ff    call   401090 <printf@plt>
4012d0: 83 fd 02          cmp    ebp,0x2
4012d3: 74 23          je    4012f8 <main+0xb8>
4012d5: bf c8 30 40 00    mov    edi,0x4030c8
4012da: e8 81 fd ff ff    call   401060 <puts@plt>
4012df: 48 81 c4 48 18 00 00    add    rsp,0x1848
4012e6: 31 c0          xor    eax,eax
4012e8: 5b              pop    rbp
4012e9: 5d              pop    rbp
4012ea: 41 5c          pop    r12
4012ec: 41 5d          pop    r13
4012ee: c3              ret
4012ef: 80 7c 24 10 00    cmp    BYTE PTR [rsp+0x10],0x0
4012f4: 74 da          je    4012d0 <main+0x90>
4012f6: eb c9          jmp   4012c1 <main+0x81>
```

As suggested by a student in class, the comparison of the user's input with the hidden password is performed just after the call to `scanf()` using the `xor` instructions:

```

40129a:    e8 11 ff ff ff      call   4011b0 <__isoc99_scnaf@plt>
40129f:    48 b8 4d 79 53 65 63  movabs rax,0x746572636553794d
4012a6:    72 65 74
4012a9:    48 33 04 24      xor    rax,QWORD PTR [rsp]
4012ad:    48 ba 50 61 73 73 77  movabs rdx,0x64726f7773736150
4012b4:    6f 72 64
4012b7:    48 33 54 24 08      xor    rdx,QWORD PTR [rsp+0x8]
4012bc:    48 09 d0      or     rax,rdx
4012bf:    74 2e      je    4012ef <main+0xaf>
...
4012ef:    80 7c 24 10 00      cmp    BYTE PTR [rsp+0x10],0x0  # <-- <main+0xaf>
4012f4:    74 da      je    4012d0 <main+0x90>

```

First, `movabs rax,0x746572636553794d` moves the hexadecimal value `0x746572636553794d` into the register `rax` (`movabs` is a synonym for `mov`). Since we are on a little-endian ISA, that value corresponds to the bytes `0x4d`, `0x79`, `0x53`, `0x65`, `0x63`, `0x72`, `0x65`, `0x74`. Referring an ASCII table, we can find that they correspond to the characters that form the first half of the password. Then, the instruction `xor rax,QWORD PTR [rsp]` does a bitwise `xor` of `rax` with the first 8 bytes at `[rsp]`, i.e. the first eight bytes of the buffer passed to `scanf`. If those bytes have a value exactly equal to that if `rax`, the result of the `xor` (stored in `rax` again) will be zero. It will be nonzero otherwise. The same thing happens for the next 8 bytes and `rdx`: they are compared to `0x64726f7773736150`, i.e. `0x50`, `0x61`, `0x73`, `0x73`, `0x77`, `0x6f`, `0x72`, `0x64`. If both `rax` and `rdx` are zero, then we jump to `main + 0xaf` and check that `[rsp+16]` is zero. Thus, our password is 16 bytes long with values: `0x4d`, `0x79`, `0x53`, `0x65`, `0x63`, `0x72`, `0x65`, `0x74`, `0x50`, `0x61`, `0x73`, `0x73`, `0x77`, `0x6f`, `0x72`, `0x64`.

- Exercise 8: I used the `-size` parameter of the `magick` command whereas I should have used the `-scale` parameter. The final `Makefile` is as follows:

```

IMAGES := $(shell ls source/*.jpeg)
THUMB := $(IMAGES:source/%.jpeg=thumbnail/%.jpeg)

.PHONY: default

default: $(THUMB)

$(THUMB) : thumbnail/%.jpeg : source/%.jpeg Makefile
    magick $(<) -scale "256x256" $(@)

```

A few links to the `magick` documentation:

- General `magick` command documentation
- `-scale` parameter
- `-size` parameter
- geometry specification